

# Code-Based Cryptography

1. Error-Correcting Codes and Cryptography
2. McEliece Cryptosystem
3. Message Attacks (ISD)
4. **Key Attacks**
5. Other Cryptographic Constructions Relying on Coding Theory

## 4. Key Attacks

1. Introduction
2. Support Splitting Algorithm
3. Distinguisher for GRS codes
4. **Attack against subcodes of GRS codes**
5. Error-Correcting Pairs
6. Attack against GRS codes
7. Attack against Reed-Muller codes
8. Attack against Algebraic Geometry codes
9. Goppa codes still resist

# Subcodes of GRS codes for the McEliece scheme



## Subcodes of GRS codes



T. Berger and P. Loidreau.

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## **Attack against this proposal:**



C. Wieschebrink.

*Cryptanalysis of the Niederreiter public key scheme based on GRS subcodes.*

In Post-Quantum Cryptography, volume 6061 of Lecture Notes in Comput. Sci., pages 61—72, 2010.

## Attack - If $2k - 1 \leq n - 2$

Public Key:  $\mathcal{K}_{\text{pub}} = \left\{ \begin{array}{l} \text{a gen. matrix of } \mathcal{C} \subseteq \text{GRS}_k(\mathbf{a}, \mathbf{b}) \\ \text{and } t = \left\lfloor \frac{n-k}{2} \right\rfloor \end{array} \right.$

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With **High Probability**:

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STEP 2 Apply the **Sidelnikov-Shestakov** attack to recover

**a**    and    **b \* b**

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## Shortened code $S_J(\mathcal{C})$

The words of  $S_J(\mathcal{C})$  are codewords of  $\mathcal{C}$  that have a zero in the  $J$ -locations, i.e.

$$S_J(\mathcal{C}) = \{\mathbf{c}_{\bar{J}} \mid \mathbf{c} \in \mathcal{C} \text{ and } c_j = 0 \text{ for all } j \in J\}$$

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$$G = \begin{array}{c|c} \begin{array}{ccc} 1 & & 0 \\ & \ddots & \\ 0 & & 1 \end{array} & \begin{array}{c} \text{ } \end{array} \\ \hline \begin{array}{ccc} 0 & \dots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \dots & 0 \end{array} & \begin{array}{c} \text{ } \end{array} \end{array}$$

Diagram illustrating the structure of the generator matrix  $G$  for a shortened code. The matrix is partitioned into four blocks. The top-left block is an identity matrix of size  $|J| \times |J|$ . The top-right block is a  $|J| \times (n - |J|)$  matrix. The bottom-left block is a  $(k - |J|) \times |J|$  matrix. The bottom-right block is a  $(k - |J|) \times (n - |J|)$  matrix. The dimensions are indicated by arrows:  $|J|$  for the top-left block,  $n - |J|$  for the top-right block,  $|J|$  for the bottom-left block, and  $k - |J|$  for the bottom-right block.

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Diagram illustrating the structure of the generator matrix  $G$  for the shortened code  $S_J(\mathcal{C})$ . The matrix is partitioned into four blocks. The top-left block is an identity matrix of size  $|J| \times |J|$ . The top-right block is a  $|J| \times (n - |J|)$  matrix. The bottom-left block is a  $(k - |J|) \times |J|$  matrix of zeros. The bottom-right block is a  $(k - |J|) \times (n - |J|)$  matrix labeled "Generator matrix for  $S_J(\mathcal{C})$ ". The total number of rows is  $k$  and the total number of columns is  $n$ .

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# Shortening a GRS code

The shortened code of a GRS code is GRS code

For GRS code we always have:

$$S_J(\text{GRS}_k(\mathbf{a}, \mathbf{b})) = \text{GRS}_{n-|J|}(\mathbf{a}_{\bar{J}}, \mathbf{b}') \text{ with } b'_i = b_i \prod_{j \in J} (a_i - a_j)$$

**Proof:** Assume  $J = \{1\}$ . Let  $G$  be a gen. matrix for  $\text{GRS}_k(\mathbf{a}, \mathbf{b})$ .

$$G = \begin{array}{cccc} b_1 & b_2 & \dots & b_n \\ b_1 a_1 & b_2 a_2 & \dots & b_n a_n \\ \vdots & \vdots & \ddots & \vdots \\ b_1 a_1^{k-1} & b_2 a_2^{k-1} & \dots & b_n a_n^{k-1} \end{array}$$

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$$\text{Thus, } g'_{ij} = \begin{cases} 0 & \text{if } j = 1 \\ \underbrace{b_j (a_j - a_1)}_{b'_j} a_j^{i-1} & \text{if } j \geq 2 \end{cases}$$

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$\rightarrow \mathbf{g}'_1 = \frac{\mathbf{g}_1}{b_1}$   
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 Generator matrix  
 for  $S_1(\text{GRS}_k(\mathbf{a}, \mathbf{b}))$

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### The Algorithm:

STEP 1 Chose a set of indices

$$J = \{i_1, \dots, i_N\} \subseteq \{1, \dots, n\} \text{ such that } 2(k - N) \leq n - 2$$

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**STEP 2** Compute a generator matrix of the shortened code  $S_J(\mathcal{C})$

$$S_J(\mathcal{C}) \subseteq \text{GRS}_{k-N}(\mathbf{a}_J, \mathbf{b}')$$

Recall that

$$\text{with } b'_i = b_i \prod_{j \in J} (a_i - a_j) \text{ for all } i \in J$$



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Note that  $2(k - N) \leq n - 2$ .

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**STEP 4** Return to **STEP 1** until  $\mathbf{a}$  is completely retrieved.

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