Code-Based Cryptography

McEliece Cryptosystem

Inria informatiques mathématiques

2. McEliece Cryptosystem

Circulant matrix

$$A = \begin{pmatrix} a_0 & a_1 & a_2 & \cdots & a_{r-1} \\ a_{r-1} & a_0 & a_1 & \cdots & a_{r-2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ a_1 & a_2 & a_3 & \cdots & a_0 \end{pmatrix}$$

Circulant matrix

$$a(X) = a_0 + a_1 X + a_2 X^2 + \ldots + a_{r-1} X^{r-1}$$
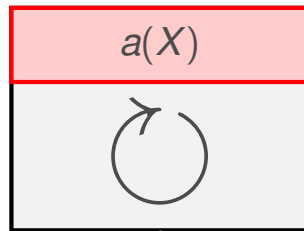
Polynomial Representation

$$A = \begin{pmatrix} a_0 & a_1 & a_2 & \cdots & a_{r-1} \\ a_{r-1} & a_0 & a_1 & \cdots & a_{r-2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ a_1 & a_2 & a_3 & \cdots & a_0 \end{pmatrix}$$

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i-th row: $X^i a(X) \mod (X^r - 1)$

1

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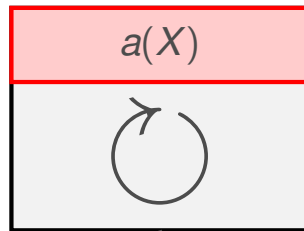
i-th row: $X^i a(X) \mod (X^r - 1)$

Proposition:
Circulant matrices of size $r \times r$ $\simeq$ Polynomials in $\mathbb{F}_q[X]/X^r - 1$

Quasi-cyclic codes

Block-Circulant Matrix

| $g_1(X)$ | $g_2(X)$ | $g_3(X)$ |
| $g_4(X)$ | $g_5(X)$ | $g_6(X)$ |

Quasi-Cyclic codes

A Linear code that admit a block-circulant parity check matrix.

Quasi-cyclic codes

1		$g_1(X)$	$g_2(X)$	$g_3(X)$
$\ddots$				
1				
	1	$g_4(X)$	$g_5(X)$	$g_6(X)$
	$\ddots$			
	1			

Quasi-Cyclic codes

A Linear code that admit a block-circulant parity check matrix.

Variants based on Algebraic codes with symmetry

➤ Using subcodes of BCH codes

P. Gaborit.
Shorter keys for code based cryptography.
In International Workshop on Coding and Cryptography (WCC 2005), pp. 81-91.

n	t	Claimed security	Public-Key sizes
2047	31	2^{80}	40505 bits
4095	26	2^{90}	12302 bits

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✗ Attack against this proposal:

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Cryptanalysis of two McEliece cryptosystems based on quasi-cyclic codes.
Special Issues of Mathematics in Computer Science, 3(2), pp. 129-140. 2010.

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Weakness of the proposal:

1. The public code comes from a primitive BCH code

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Weakness of the proposal:

1. The public code comes from a primitive BCH code
2. The permutation (used to hide the secret code) is too restrictive

Variants based on Algebraic codes with symmetry

➤ Using subcodes of QC-Alternant codes

📄 T.P. Berger, P.L. Cayrel, P. Gaborit and A. Otmani.
Reducing key length of the McEliece cryptosystem.
In AFRICACRYPT 2009, pp. 77-97.

q	n	k	t	Security	Public-Key sizes
2^8	663	561	25	2^{80}	8980 bits
2^8	663	510	37	2^{95}	12240 bits
2^8	1020	867	37	2^{116}	20800 bits

✗ Attack against this proposal:

📄 J.C. Faugère, A. Otmani, L. Perret and J.P. Tillich.
Algebraic cryptanalysis of McEliece variants with compact keys.
In EUROCRYPT 2010, pp. 279-298.

Variants based on Algebraic codes with symmetry

Idea of the attack:

Solve the system

$$G\, H^T = 0$$

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Public generator matrix

Variants based on Algebraic codes with symmetry

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$$G\ H^T \leftarrow 0$$

Public generator matrix

Unknown alternant parity check matrix

$$H = \begin{pmatrix} 1 & 1 & \dots & 1 \\ a_1 & a_2 & \dots & a_n \\ a_1^2 & a_2^2 & \dots & a_n^2 \\ \vdots & \vdots & \ddots & \vdots \\ a_1^{r-1} & a_2^{r-1} & \dots & a_n^{r-1} \end{pmatrix} \begin{pmatrix} b_1 & & & 0 \\ & b_2 & & \\ & & \ddots & \\ 0 & & & b_n \end{pmatrix}$$

Variants based on Algebraic codes with symmetry

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The quasi-cyclic structure allows a drastic reduction of the number of unknowns

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